# How They Do It: Unix Hacking 101

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## Who Am I?

- Independent security consultant
- SANS Institute Senior Faculty
- Technical Editor for Sys Admin
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Generally speaking, a guy who probably spends way too much time with Unix...

#### What's In This Course?

- Overview of common techniques for breaking into Unix systems
- Detailed coverage of stack smashing (buffer overflow) attacks and defenses
- A look at what happens after a successful break-in: rootkit installation
- Live demos, prevention techniques, etc.

## What's Your Job?

# ASK QUESTIONS!

# **Overview**

# Hierarchy of Vulnerabilities

- Physical access
- Captured (or weak/default) password
- Deliberate malware
- Software vulnerabilities (including race conditions, buffer overflows, et al)
- Subverted "trust relationships" (.rhosts files, X Windows authentication, etc.)
- Session hijacking

# Physical Access = root

- Single-user boot
- Boot off of OS Media
- Corrupt root file system
- Steal the disk drives!

Keep those critical servers locked up!

# Corrupt Root File System

- Repeatedly power-cycle system
- Root file system will eventually become inconsistent— manual fsck required
- System will come up at root shell without asking for a password
- Attacker can fsck file system and change root password, etc.

#### Weak Passwords

- We know reusable passwords are bad
- Better authentication systems are expensive to implement
- Common solutions are far from ideal:
  - Account expiration, password change
  - "Cracking" passwords and harassing users
- Unfortunately, many Unix systems lack ability to force "strong" passwords

## Deliberate Malware

- Common Unix vectors:
  - Trojan software distributions (often after break-ins at software archive servers)
  - Trap-doors in web programming libraries
- Critical to verify software signature (with PGP, if possible) before installing
- Consider testing on isolated network before production deployment

# File System Race Conditions

- Typically caused by programs writing to directories with "unsafe" permissions
- Compounded by programs which choose "predictable" file names
- Attacker creates a symlink which causes program to modify an unexpected file
- Window of vulnerability can be very small
   perhaps a few CPU cycles

## 1992 SunOS /bin/mail Hole

- | /bin/mail called by Sendmail to
  | deliver mail into /var/spool/mail
- /bin/mail checks mail file to see if it's a symlink and then opens file
- Between the check and the open, attacker creates a symlink to / .rhosts

## Other Predictable File Names

- mktemp() generates temporary file names based on template string
- Attacker can usually predict next file name
- Attacker makes link after mktemp() call and before open()
- Use mkstemp() instead or at least call
  open() with O CREAT | O EXCL set

# Set-UID Script Race Condition

- Make a symlink to set-UID script
- Execute symlink
- While kernel is loading interpreter, re-point symlink to root shell script
- New script executes with set-UID permissions from old script!

## Trust Problems - . rhosts

- This mechanism replaces password authentication with IP/host-based auth:
  - IP addresses can be spoofed
  - DNS can be corrupted
  - Root privs on other hosts can be exploited
- Attacker can often "reverse" .rhosts files and find other machines to break

# X Windows Exploits

- Attacker controlling your display can:
  - Get remote video output
  - Capture all keystrokes (read passwords)
- "All or nothing" access model is a real problem here

# Session Hijacking

- Attacker "takes over" session in progress:
  - Attacker doesn't have to guess passwords
  - Pretty obvious to affected user, however
- Two types of session hijacking attacks:
  - Local attacks snoop streams via kernel
  - Remote attacks from third-party machines

# Protecting Yourself

- Patching is important, but it only protects you from known vulnerabilities
- Disabling services you're not using protects you from as yet unknown attacks
- Use firewalls at the network and host level to control access and drop bogus traffic
- Encrypt all network traffic with SSH, SSL, IPSEC, or other strong VPN

## Extra Credit Items

- Use a strong, two-factor authentication system for user access (expensive)
- Abandon .rhosts files in favor of DSA authentication via SSH (requires re-tooling)
- Monitor system configurations with a tool like AIDE or Tripwire™ (mgmt issues)

# Stack Smashing

# The History of Buffer Overflows

- First wide-spread buffer overflow attack was the 1988 Morris Worm (fingerd)
- Mudge's white paper in 1995 apparently popularized the term "buffer overflow"
- Buffer overflow exploits on Linux/Solaris used to motivate Y2K DDoS attacks
- Fully automated worms now automatically infecting systems via these vulnerabilities

# **Process Memory**

Text

Data

BSS

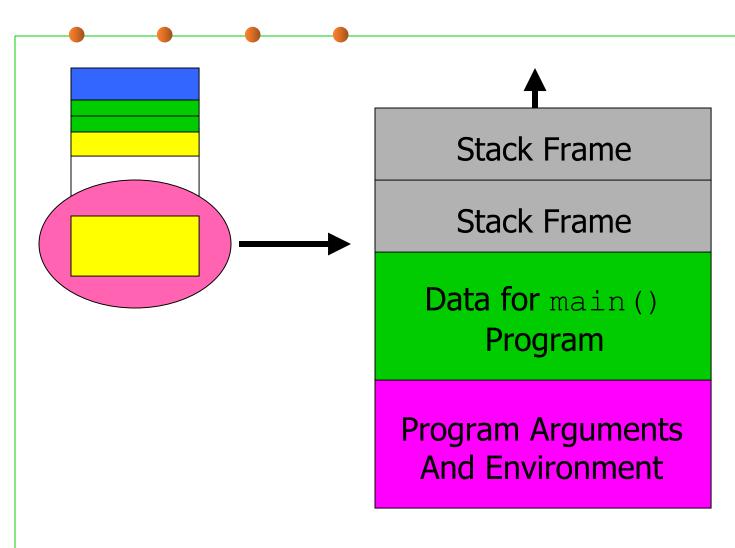
Heap



Stack

- ← Program instructions (read-only)
- ← "Constants" literal strings/numbers
- ← Global and persistent (static) variables
- ← Dynamically allocated data
- ← Free/unallocated memory
- ← Local subroutine data

## The Stack



One frame added per subroutine call

## Stack Frame

Instruction to return to when subroutine exits

Subroutine Variables

Frame Pointer

Return Address

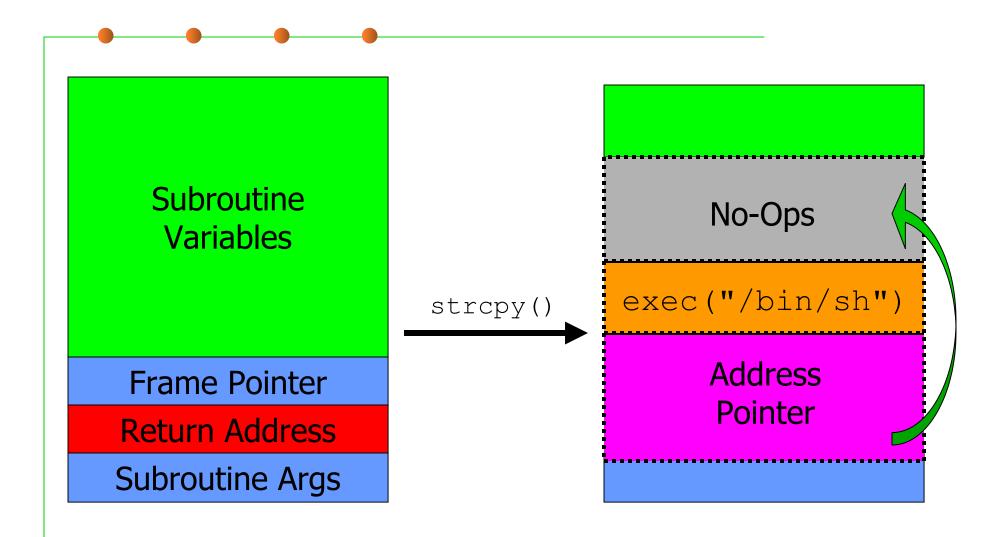
Subroutine Args

Data copied into buffers from top of frame to bottom

## Classic Buffer Overflow

- Attacker constructs a string containing:
  - No-ops for padding
  - Machine code to exec("/bin/sh")
  - Bogus instruction addr pointing into subroutine data area
- Subroutine is coerced into copying this string into its data area, overwriting end of buffer
- Subroutine exits and program follows bogus address to execute shell
- If program runs as root, attacker gets root shell

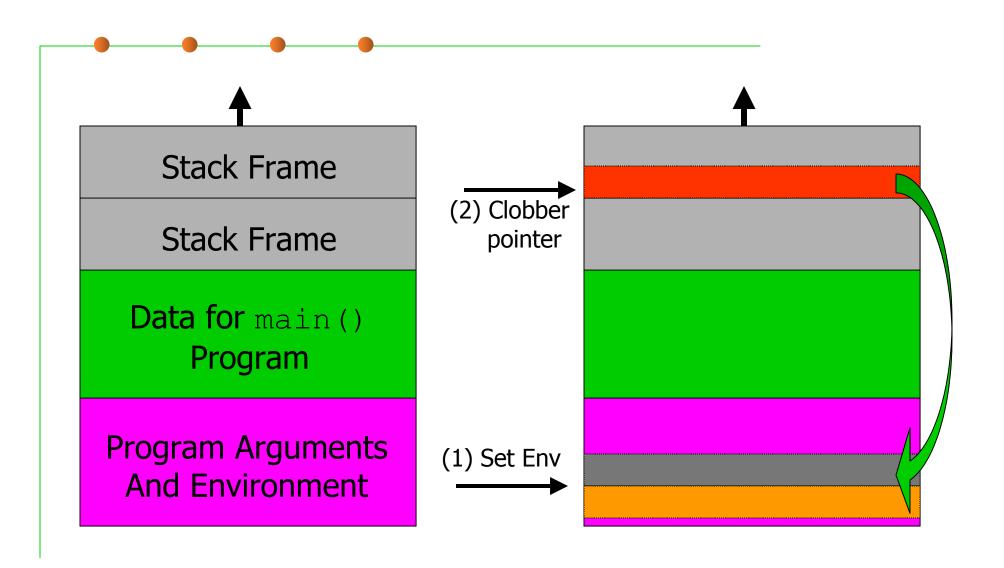
# Classic Buffer Overflow (2)



## **Modified Attack**

- Sometimes subroutine buffer is too small to hold exploit code
- Exploit code can be put into an environment variable
- Subroutine return address set to environment area at bottom of stack
- Requires attacker have local machine access...

# Modified Attack (2)



## In General...

- Attacker places code in program memory, changes address pointer to jump to code
- Places to put code:
  - Stack (per previous examples)
  - Heap
  - Data/BSS (if writable)
- Ways to modify address pointer:
  - Stack overflow (per previous examples)
  - Format string attacks (see next slide)

# Format String Attacks

Programmers are sometimes lazy:

```
printf("%s", str);  # correct
printf(str);  # lazy
```

- If attacker can set str, then attacker can embed output specifiers
- In particular, %n writes a numeric value to a specified memory address
- Attacker can overwrite return address pointer or other numeric value

# Fixing Buffer Overflows

- Fix the programs
  - Too many
  - New bugs being written daily
- Fix the programmers
  - Too many
  - New programmers being made daily
- Fix the stack

# Fixing the Stack

- Code normally resides in read-only text area at beginning of program memory
- Well-behaved programs should never execute instructions off data stack
- Modify kernel— attempts to execute from stack cause program to abort
- Requires cooperation of CPU hardware

# Different Implementations

- For Solaris, HP-UX, et al:
  - "Stack protection" only prevents executing code off of stack pages
  - Can be thwarted by putting malicious code into the heap area
- OpenBSD (W^X) and Adamantix (PaX):
  - All writable pages marked as non-executable (including heap area)
  - Can break applications written in high-level languages (notably Java)

# Compiler-Based Solutions

- Compiler inserts a random "canary" value before address pointer
- Canary value checked when subroutine exits – abort if changed
- Attacks which clobber return addr pointer also clobber canary
- Performance hit due to extra code to insert/check canary values

Subroutine Variables

Frame Pointer

"Canary"

Return Address

Subroutine Args

#### How to Defeat Stack Protection

- Put malicious code into heap data area (at least for Solaris and HP-UX)
- Force subroutine return to other call in text segment, like system()
- Overwrite expected subroutine args:
  - system()/exec() call wrong program
  - open () opens unexpected file

## After the Attack

#### **Back Doors**

- Having broken a system, the attacker wants to make it easy to get back in
- Often sshd or telnetd replaced, but any networked or SUID binary will do
- New version gives root shell when special account and/or password used
- Multiple back doors usually left behind to "fake out" system admins

#### rootkits

- Nowadays, attackers want to install IRC servers, sniffers, DDOS tools, et al
- Need much more intricate tools to cover their tracks
- rootkits are pre-packaged bundles of software with tools and Trojans
- New trend is rootkits that "harden" the system and remove other rootkits!

# Hiding Files

- ls, dir, find, etc. all replaced to hide
  attackers' files
- Some rootkits use configuration files to change behavior "on-the-fly"
- Binaries are crafted to have the same size (and same timestamps) as originals
- Some files are installed "immutable" to make them harder to remove

# Hiding Processes

- ps, top, lsof commands replaced with versions to hide certain processes
- Also need to modify kill, pkill, killall and similar programs

# **Hiding Network Connections**

- New syslogd and tcpd will not log connections from certain hosts/domains
- ifconfig will not report that network interfaces are in promiscuous mode
- netstat will not show connections from certain hosts/domains

### Problems for Rootkit Authors

- Too many binaries to replace:
  - Hard work
  - Greater chance of detection
  - Might miss one
- Checksumming (with AIDE or Tripwire) allows for easy detection

### The Next Wave – Kernel Hacks

- Suppose you could replace the underlying kernel routines?
- Single modification corrupts every program on the system!
- Modern Unix systems support dynamic kernel modules— no reboot required!
- New kernel rootkits can even attack via /dev/kmem- no loadable modules!

### What Do You Do?

- Analyze system by first booting from OS CD-ROM
- All tools should be brought over on read-only media (e.g., CD-ROM)
- Reinstall system from scratch don't necessarily trust your backups

# And Don't Forget...

- Assume any passwords on compromised system have been cracked off-line
- If you believe packet sniffer was installed, then <u>all</u> site passwords must change

# Wrap Up

### That's All Folks!

- Any final questions/comments?
- Please fill out your eval forms!
- Thanks for listening!

Plenty of useful URLs to follow...

#### General Resources

- SANS "Reading Room" http://www.sans.org/rr/
- CERT/CC "Tech Tips"
  http://www.cert.org/tech\_tips/
- Center for Internet Security http://www.CISecurity.org/
- BUGTRAQ Mailing List Archives http://www.securityfocus.com/

# Details on Stack Smashing

- The classic paper by Aleph1: http://www.phrack.org/show.php?p=49&a=14
- Format string attacks http://www.securityfocus.com/guest/3342
- Heap and BSS overflow paper: http://w00w00.org/files/articles/heaptut.txt
- libc redirection attacks:
  http://hackersplayground.org/papers/stack.txt

#### Other Resources

- Layered defense from OpenBSD 3.[34]: http://www.openbsd.org/papers/csw03.mgp
- Compiler-based stack protection:
   StackGuard http://www.immunix.org/
   IBM http://www.trl.ibm.com/projects/security/ssp/
- Session hijacking tools:
   Hunt http://lin.fsid.cvut.cz/~kra/#HUNT
   Ettercap http://ettercap.sourceforge.net/

#### **Rootkit Information**

Links to detailed rootkit overviews:

http://www.chkrootkit.com/

http://biocserver.bioc.cwru.edu/~jose/shaft\_analysis/node-analysis.txt

Functional kernel rootkits (Linux):

http://la-samhna.de/library/rootkits/list.html

http://www.phrack.org/show.php?p=58&a=7